Fachbereich

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## Assignment 2 - Bonus Material Cryptography

In the following you will find some bonus material to the second exercise including Modular Exponentiation and Fermat's factorization.

## Modular Exponentiation

The fast exponentiation algorithm is based on the assumption that for

$$
x^{a} \bmod n=\left(\left(\left(x^{2} \bmod n\right)^{2} \bmod n\right)^{2} \ldots\right)^{2} \bmod n, \quad \text { with } a=2^{k}(a \text { power of } 2)
$$

squaring is iterated k times. For example, $x^{\left(2^{2048}\right)}$ can be computed in 2048 modular squaring operations (Knospe 2019, p. 67).

For exponents that are not a power of 2, we can write it as a sum of powers of 2. (Remember that $\left.x^{b} \cdot x^{c}=x^{b+c}\right)$

## Example (Lecture Slide 21):

To calculate $7^{17} \bmod 77$ we can write $17=2^{4}+2^{0}$.
We also have $2^{4}=2^{2} \cdot 2^{2}$ and $2^{2}=2^{1} \cdot 2^{1}$.

| Powers of 2 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 16 | 8 | 4 | 2 | 1 |
| $2^{4}$ | $2^{3}$ | $2^{2}$ | $2^{1}$ | $2^{0}$ |

$7^{1} \equiv 7 \bmod 77 \Leftrightarrow 7 \bmod 77=7$
$7^{2} \equiv 49 \bmod 77 \Leftrightarrow 49 \bmod 77=49$
$7^{4} \equiv 7^{2} \cdot 7^{2} \equiv 49 \cdot 49 \bmod 77 \Leftrightarrow 2401 \bmod 77=14$
$7^{8} \equiv 7^{4} \cdot 7^{4} \equiv 14 \cdot 14 \bmod 77 \Leftrightarrow 196 \bmod 77=42$
$7^{16} \equiv 7^{8} \cdot 7^{8} \equiv 42 \cdot 42 \bmod 77 \Leftrightarrow 1764 \bmod 77=70$
$7^{17} \equiv 7^{16} \cdot 7^{1} \equiv 70 \cdot 7 \bmod 77 \Leftrightarrow 490 \bmod 77=28$

## Bonus Exercise:

Use modular exponentiation to solve $4^{64} \bmod 7$ and $3^{210} \bmod 15$. You can check your results on https://sagecell.sagemath.org with the following syntax (e.g. $3^{128} \bmod 15$ ):

Type some Sage code below and press Evaluate.


## 6

Help I Powered by SageMath

## Fermat's Factorization

Let $p, q$ be prime and $N=p q$. Fermat's factoring represents $N$ as a difference of 2 squares:

$$
\left.N=x^{2}-y^{2}=(x+y)(x-y)\right)
$$

First, we start with $\mathrm{x}=\left\lceil\sqrt{N} \mid\right.$ and then increase x by 1 until $x^{2}-N$ is square (so that we can derive y) so that $\boldsymbol{N}=x^{2}-y^{2}$ holds.

This method works because we can represent $N$ as a difference of 2 squares:

$$
p q=\left(\frac{1}{2}(p+q)^{2}-\left(\frac{1}{2}(p-q)^{2}\right)^{2}=x^{2}-y^{2} .\right.
$$

You will find this explanation with more details in Knospe 2019, p. 178 f.

## Bonus Exercise:

Let $N=247$; then we first set $x \approx \sqrt{N}$. We obtain $x=16$ and derive $x^{2}-N=9$. Because 9 is square we know that $y=3$. From above we know that $p q=(x+y)(x-y)$ so we receive $247=$ 19-13.

If you need more practice, just choose some primes $p, q$ and derive $N$. Then use Fermat's factoring to solve $N$.

## Literature

(Knospe 2019) Knospe, Heiko. A Course in Cryptography. Vol. 40. American Mathematical Soc., 2019. https://ebookcentral.proquest.com/lib/senc/reader.action?docID=5962876

