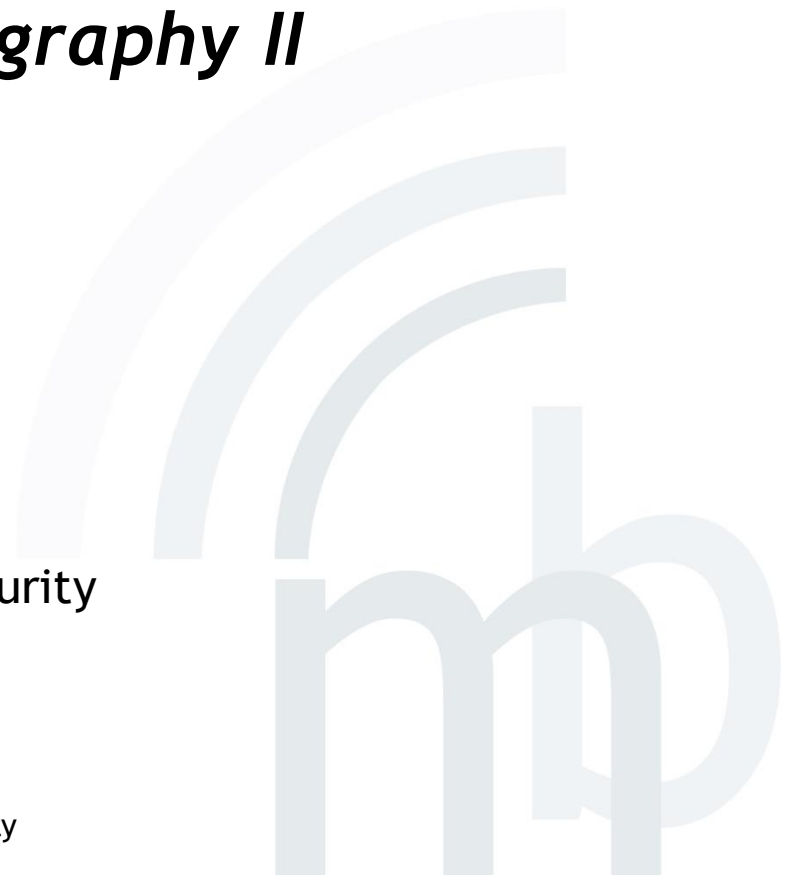


# *Assignment 4 - Cryptography II*

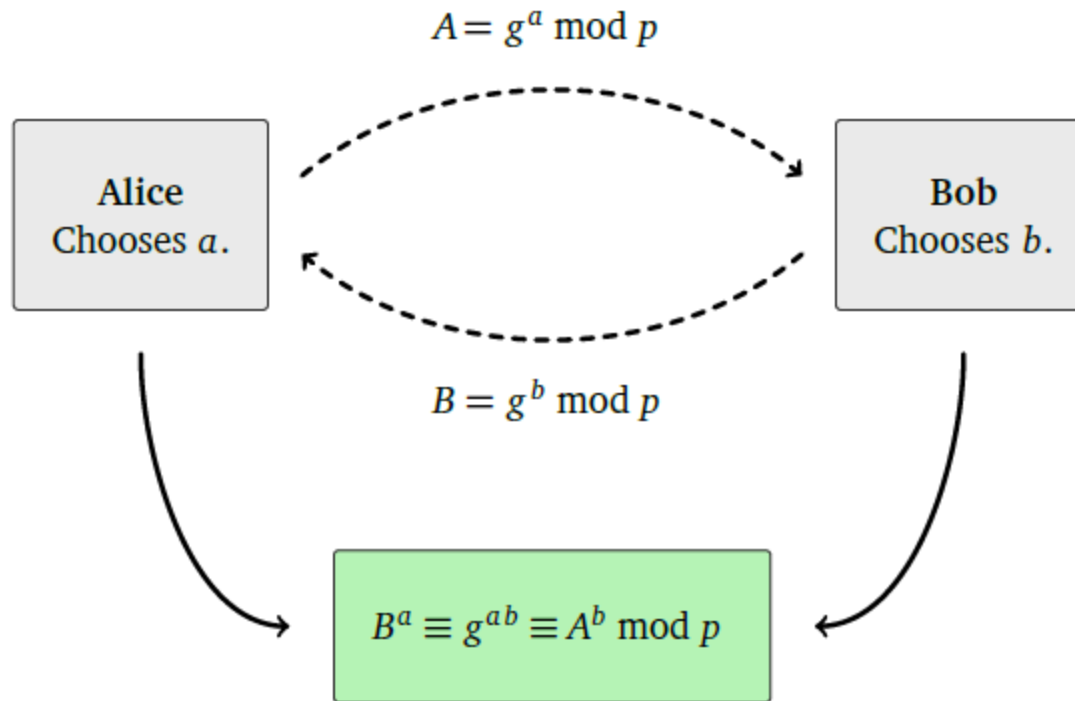
Information & Communication Security  
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Describe Diffie-Hellman key exchange method in detail.



Diffie-Hellman Key Exchange

## Exercise 1 - Solution (Cont.)

- a) Alice and Bob agree publicly on a cyclic group, e.g.  $G = \langle g \rangle$ ,  $G = \mathbb{F}^*$ .
- b) Alice chooses randomly some  $0 \leq a < |G|$  and computes  $A := g^a$ . Bob chooses randomly some  $0 \leq b < |G|$  and computes  $B := g^b$ .
- c) Alice sends Bob  $A$ . Bob sends Alice  $B$ .
- d) Alice computes  $S := B^a = (g^b)^a = g^{ab}$ . Bob computes  $S := A^b = (g^a)^b = g^{ab}$ .
- e) Now Alice and Bob can use  $S$  as their secret key to encrypt and decrypt messages.

## Exercise 1 - Solution (Cont.)

Outside of this process Eve only knows  $G = \langle g \rangle$ ,  $A$  and  $B$ , but she does not know  $a$ ,  $b$ ,  $S$ . Thus Eve either needs to compute  $a = \log_g A$  and  $b = \log_g B$  (this is known as *discrete logarithm problem* and is assumed to be “hard”); or she has some other magical function  $f$  such that  $S = f(A, B, G)$ . Clearly, security of this system highly relies on the choice of the group, i.e.  $g$ . For example, taking  $G = (\mathbb{Z}/n\mathbb{Z}, +) = \langle 1 \rangle$ , thus exponentiation  $g^a$  boils down to  $g \cdot a = a$  in this setting.

In order to prepare to receive encrypted messages with the RSA cryptosystem, Alice has chosen primes  $p = 23$  and  $q = 37$ . She has also chosen  $e = 13$  as her public key (also called her public exponent).

a) Determine Alice's public modulus  $n$ .

- <https://www.khanacademy.org/computing/computer-science/cryptography/modern-crypt/v/intro-to-rsa-encryption>
- <http://mathworld.wolfram.com/RSAEncryption.html>
- [https://simple.wikipedia.org/wiki/RSA\\_algorithm](https://simple.wikipedia.org/wiki/RSA_algorithm)

## Exercise 2 - Solution (a)

Alice's public modulus is  $n = p \cdot q = 23 \cdot 37 = 851$ .



b) Suppose that Bob wants to send Alice the message BAT. Determine the base twenty-six representation of the ciphertext that he will send to Alice.

## Exercise 2 - Solution (b)

Converting the message BAT to base twenty-six gives  $BAT = 1 \cdot 26^2 + 0 \cdot 26 + 19 = 695$ .

Therefore, since the RSA encryption function is  $c = m^e \text{ MOD } n$ , Bob's ciphertext is given by  $c = 695^{13} \text{ MOD } 851$ . In order to compute the modular exponential, we repeatedly square:

$$\begin{aligned} 695^2 &\equiv 508 \pmod{851} \\ 695^4 &\equiv 508^2 \equiv 211 \pmod{851} \\ 695^8 &\equiv 211^2 \equiv 269 \pmod{851}. \end{aligned}$$

We then conclude that

$$695^{13} \equiv 695^8 \cdot 695^4 \cdot 695 \equiv 269 \cdot 211 \cdot 695 \equiv 593 \cdot 695 \equiv 251 \pmod{851}.$$

Thus, the ciphertext is  $c = 251$  which converted to base twenty-six reads

$$251 = 9 \cdot 26 + 17 = \text{JR}.$$

c) Determine Alice's private key (or decryption key) d.

## Exercise 2 - Solution (c)

Alice's decryption key is given by  $d = e^{-1} \text{MOD} \varphi(n)$  where  $\varphi(n) = (p - 1) \cdot (q - 1)$ . We find  $\varphi(n) = 22 \cdot 36 = 792$ , so that  $d = 13^{-1} \text{MOD} 792$ . In order to calculate this modular inverse, we use the extended Euclidean algorithm so that:

$$\begin{aligned} 792 &= 60 \cdot 13 + 12 \\ 13 &= 12 + 1. \end{aligned}$$

Back substitution therefore gives:

$$-792 + 61 \cdot 13 = 1$$

from which we conclude that

$$d = 13^{-1} \text{MOD} 792 = 61.$$

d) Suppose that Bob has also sent Alice the ciphertext  $y = 625$ . Determine the base twenty-six representation of the plaintext message.

## Exercise 2 - Solution (d)

The RSA decryption function for Alice is given by  $m = c^d \text{ MOD } n$  so that Bob's plaintext message is  $m = 625^{61} \text{ MOD } 851$ . In order to compute the modular exponential, we repeatedly square:

$$625^2 \equiv 16 \pmod{851}$$

$$625^{16} \equiv 9^2 \equiv 81 \pmod{851}$$

$$625^4 \equiv 16^2 \equiv 256 \pmod{851}$$

$$625^{32} \equiv 81^2 \equiv 604 \pmod{851}.$$

$$625^8 \equiv 256^2 \equiv 9 \pmod{851}$$

We then conclude that

$$\begin{aligned} 625^{61} &\equiv 625^{32} \cdot 625^{16} \cdot 625^8 \cdot 625^4 \cdot 625 \equiv 604 \cdot 81 \cdot 9 \cdot 256 \cdot 625 \\ &\equiv 784 \pmod{851}. \end{aligned}$$

Thus, the plaintext is  $x = 784$  which converted to base twenty-six reads

$$784 = 1 \cdot 26^2 + 4 \cdot 26 + 4 = \text{BEE}.$$

Install PGP Email Desktop (trial version) or a similar software for mail encryption on your system. Create a new key pair, and send a signed and encrypted message to [abtin.shahkarami@m-chair.de](mailto:abtin.shahkarami@m-chair.de) containing your newly created public key and a short summary of your experiences.

PGP can be downloaded from

<http://www.symantec.com/business/desktop-email>.

Looking forward to your Emails 😊

Thank you!

- Questions: [security@m-chair.de](mailto:security@m-chair.de)